Optimizing Hash Partitioning for Solid State Drives

in heol Shin Hong han oh onmoo **Dept** omp ter S ien e **Dept** of omp ter S ien e οf le troni s onsei niversit onsei niversit ang ae daero gil Seodaem n Seodaem Seogho onsei ro onsei Seo I ep i li Seo I orea Seo I orea orea fallsmal smanioso onsei S onsei ommoo nq lge om Sangh n Par **Dept** of omp ter S ien e onsei niversit Seodaem n onsei ro Seo I orea sangh n S onsei

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The use of flashSSDs has increased rapidly in a wide range of areas due to their superior energy efficiency, shorter access and higher and idth hen ompared to HDDs he internal parallelism created by multiple flash memory packages embedded in a flashSSDs, is one of the unique features of flashSSDs. Many new DBMS technologies have been developed for flashSSDs, but query processing for flashSSDs have dra n less attention than other D S te hnologies Hash partitioning is pop larl sed in er pro essing al gorithms to materialize their intermediate results in an effiient manner n this paper e propose a novel hash par titioning algorithm that e ploits the internal parallelism of flashSSDs. The devised hash partitioning method outperforms the traditional hash partitioning te hni e regardless of the amo nt of availa le main memor independent! from the buffer management strategies (blocked I/O vs page sized

O e implemented o r method ased on the so r e ode of the PostgreS storage manager PostgreS relation as files created by the TPC-H workload were employed in the e periments O r method as fo nd to e p to tifaster than the traditional method ith lo ed O and times faster than the traditional method ith page.

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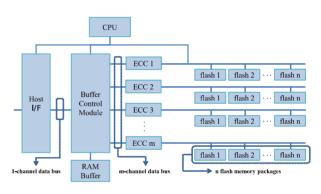
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in a ide range of areas e a se of their s perior feat res hen ompared to HDDs Personal sers have emplo ed flashSSDs in their laptop and desktop computers in order to oost operating s stem and speed p a essto fre entl used applications. Data centers have also adopted flashSSDs to improve energy efficiency and performance. Many DBMS technologies have been developed for flashSSDs. While majority of previous studies have focused on the buffer and index management of DBMSs on flashSSDs, query processing algorithms for flashSSDs have drawn less attention from resear hers

he nternal Parallelism of lashSSDs

A typical flashSSD architecture is shown in Figure 1. A flashSSD includes a CPU, a host interface (host I/F), a RAM buffer, Error-Correcting Code(ECC) modules, multiple data transfer channels and multiple flash memory pack-

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